INTEGER CHEBYSHEV POLYNOMIALS

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1. Introduction.

The ubiquitous Chebyshev polynomial

$$T_n(x) := \cos(n \arccos x)$$

$$= \frac{1}{2} \left[\left(x + \sqrt{x^2 - 1} \right)^n + \left(x + \sqrt{x^2 - 1} \right)^n \right]$$

is a polynomial of degree n with integer coefficients and with lead coefficient 2^{n-1} that equioscillates n+1 times on the interval [-1,1].

For fairly simple reasons, based on this equioscillation, it follows that

$$\min_{p_{n-1} \in \mathcal{P}_{n-1}} \|x^n - p_{n-1}\|_{[-1,1]}$$

$$= \left\| 2^{1-n} T_n \right\|_{[-1,1]} = \frac{1}{2^{n-1}}.$$

Here \mathcal{P}_n denotes the polynomials of degree at most n with real coefficient.

The polynomial

$$p_n(x) := 2\left(\frac{b-a}{4}\right)^n T_n\left(\frac{2x-a-b}{b-a}\right)$$

is now the monic polynomial of degree n of **small-est supremum norm** on the interval [a, b] and it satisfies

$$||p_n||_{[a,b]} = 2\left(\frac{b-a}{4}\right)^n.$$

So

$$(\|p_n\|_{[a,b]})^{1/n} \to \left(\frac{b-a}{4}\right)$$

which gives the **transfinite diameter** of [a, b].

The Chebyshev polynomials have a central role to play in minimization problems in the sup norm as well as many other extremal problems.

The analogue problem where the polynomials are restricted to have integer coefficients is very much harder.

We let \mathcal{Z}_n denote the polynomials of degree at most n with integer coefficients.

We define

(*)
$$\Omega_n[a,b] := \left(\inf_{0 \neq p \in \mathcal{Z}_n} \|p\|_{[a,b]}\right)^{1/n}$$

and let

$$\Omega[a, b] := \inf \{ \Omega_n[a, b] : n = 0, 1, \dots \}$$
$$= \lim_{n \to \infty} \Omega_n[a, b].$$

Any polynomial satisfying (*) above is called an n-th **integer Chebyshev polynomial** on [a, b].

The above limit exists and equals the inf mostly because

$$\left(\Omega_{n+m}[a,b]\right)^{n+m} \le \left(\Omega_n[a,b]\right)^n \left(\Omega_m[a,b]\right)^m.$$

We have from the unrestricted case the trivial inequality

$$\Omega[a,b] \ge \frac{b-a}{4}.$$

We also have

$$\Omega[a,b] \le \Omega_n[a,b]$$

for any particular n.

Thus good upper bounds can be achieved by computation (although the computation to any degree of accuracy is hard).

The limit $\Omega[a, b]$ may be thought an integer version of the **transfinite diameter**.

Hilbert showed that there exists an absolute constant c so that

$$\inf_{0 \neq p \in \mathcal{Z}_n} \|p\|_{L_2[a,b]} \le c n^{1/2} \left(\frac{b-a}{4}\right)^{1/2}$$

and Fekete showed that

$$(\Omega_n[a,b])^n \le 2^{1-2^{-n-1}}(n-1)\left(\frac{b-a}{4}\right)^{n/2}$$

There are many refinements.

From the above it follows that

$$\frac{b-a}{4} \le \Omega[a,b] \le \left(\frac{b-a}{4}\right)^{1/2}$$

Recall that $b - a \leq 4$.

There is a pretty argument due to Gelfond to see that integer coefficients really are a restriction on [0,1].

If $0 \neq p_n \in \mathcal{Z}_n$ then

$$||p_n||_{[0,1]}^2 \ge ||p_n||_{L_2[0,1]}^2 = \int_0^1 p_n^2(x) dx$$
$$= \frac{m}{\text{LCM}(1, 2, \dots, 2n+1)} \ne 0$$

where LCM denotes the least common multiple. Now LCM(1, 2, ..., n))^{1/n} $\sim e$, by the prime number theorem and it follows that

$$\Omega[0,1] \ge 1/e.$$

This is not however the right lower bound.

The best previous bounds, due to Aparicio and Gorshkov,

$$\frac{1}{(2.37686\dots)} \le \Omega[0,1] \le \frac{1}{(2.343\dots)}.$$

The upper bound comes by example.

For the lower bound we show there exist infinitely many relatively prime polynomials $q_k \in \mathcal{Z}_k$ with all their roots in (0,1), and with lead coefficients $\{a_k\}$ satisfying $a_k^{1/k} \leq 2.37686...$

This number was conjectured to be the right bound. It comes from iterating (x - 1/x) on $(-\infty, \infty)$.

Lemma. Suppose $p_n \in \mathcal{Z}_n$ and suppose $q_k(z) := a_k z^k + \cdots + a_0 \in \mathcal{Z}_k$ has all its roots in [a, b]. If p_n and q_k do not have common factors then

$$(\|p_n\|_{[a,b]})^{1/n} \ge |a_k|^{-1/k}.$$

Proof. Let $\beta_1, \beta_2, \ldots, \beta_k$ be the roots of q_k . Then

$$|a_k|^n p_n(\beta_1) p_n(\beta_2) \cdots p_n(\beta_k)$$

is a non-zero integer and the result follows. \square

Aparicio also shows that if [a, b] = [0, 1] then any polynomial $p \in \mathcal{Z}_n$ for which the inf in (*) is achieved, for sufficiently large n, has a factor of the form

$$(x)^{\lfloor \lambda, n \rfloor} (1-x)^{\lfloor \lambda_1 n \rfloor} (2x-1)^{\lfloor \lambda_2 n \rfloor}$$
$$*(5x^2 - 5x + 1)^{\lfloor \lambda_3 n \rfloor}$$

where $\lambda_1 \geq .014$, $\lambda_2 \geq .016$, $\lambda_3 \geq .0037$.

We improve the upper bounds of $\Omega[0,1]$ to

$$\Omega[0,1] \le \frac{1}{2.360}$$

and use this to increase the number of factors that must divide an n-th integer Chebyshev polynomial T_n on [0,1].

We also establish a lower bound for the multiplicity of the zero at 0 of the integer Chebyshev polynomial $T_{n,a}$ on [a,1].

From all this we deduce that the natural conjecture, due to the Chudnovsky's, that

$$\Omega[0,1] = \frac{1}{2.37686..}$$

is close but false.

These high order zeros are also used to show that the function $\Omega(x) := \Omega[0, x]$ is constant on the interval $[1 - \delta, 1 + \delta]$.

We also relate the integer Chebyshev problem on small intervals [0, 1/m] to an old problem of Schur and Siegel on the trace of totally positive algebraic integers.

Another related problem is the Prouhet-Tarry-Escott problem.

We conclude with a number of open problems.

2. Computing Integer Chebyshev Polys.

We restrict our attention to the interval [0, 1]. Though we observe in passing that

$$(\Omega[-1,1])^4 = (\Omega[0,1])^2 = \Omega[0,1/4]$$

as a consequence of the changes of variable $x \to x^2$ and $x \to x(1-x)$ and symmetry.

The dependence of the constant $\Omega[a, b]$ and the minimal polynomials on [a, b] is interesting and is explored a little further later.

Even computing low degree examples is complicated. There is no good algorithm and getting examples of say degree 100 seems intractable.

Note that we do not have uniqueness, though it is open as to whether we have uniqueness for n sufficiently large.

The arguments for the above table are of the following variety.

Consider the case n = 5.

Let $T_5 \in \mathcal{Z}_5$ be a 5-th integer Chebyshev polynomial on [0,1].

Then $T_5(0)$ and $T_5(1)$ are integers of modulus less then 1 so both of them must be 0.

Using Markov's inequality, we obtain that

$$||T_5'||_{[0,1]} \le 50||T_5||_{[0,1]} \le \frac{50}{(2.236\dots)^5} < 1.$$

Since $T_5'(0)$ and $T_5'(1)$ are integers of modulus less than 1, both of them must be zero.

Since $2^5T_5(1/2)$ is an integer of modulus at most $32/(2.236...)^5 < 1$, it must also be zero.

Thus we conclude that

$$x^{2}(1-x)^{2}(2x-1)$$
 divides T_{5} .

Examples in $L_2[0,1]$. Polynomials of degrees 13..., 20, which minimize $||p_n||_{L_2[0,1]}$ This computation was done in pari by using the minum function.

Degree 13

$$(5z^2 - 5z + 1)(2z - 1)^2z^4(z - 1)^5$$

$$(5z^2 - 5z + 1)(2z - 1)^2(z - 1)^4z^5$$

$$(5z^2 - 5z + 1)(2z - 1)^3(z - 1)^4z^4$$

$$(2z-1)(5z^2-5z+1)(z-1)^5z^5$$

$$(2z-1)(5z^2-5z+1)^2(z-1)^4z^4$$

$$(2z-1)^3(z-1)^5z^5$$

$$(2z-1)(z-1)^4z^4(29z^4-58z^3+40z^2-11z+1)$$

Degree 14

$$(5z^2 - 5z + 1)(2z - 1)^2(z - 1)^5z^5$$

Degree 15

$$(5z^2 - 5z + 1)(2z - 1)^3(z - 1)^5z^5$$

Degree 16

$$(5z^2 - 5z + 1)(2z - 1)^2(z - 1)^6z^6$$

Degree 17

$$(5z^2 - 5z + 1)(2z - 1)^3(z - 1)^6z^6$$

Degree 18

$$(2z-1)^2(5z^2-5z+1)(z-1)^6z^6$$

Degree 19

$$(5z^2 - 5z + 1)(2z - 1)^3(z - 1)^7z^7$$

Degree 20

$$(5z^2 - 5z + 1)(29z^4 - 58z^3 + 40z^2 - 11z + 1)(2z - 1)^2(z - 1)^6z^6$$

Let

$$p_0(x) := x$$

$$p_1(x) := 1 - x$$

$$p_2(x) := 2x - 1$$

$$p_3(x) := 5x^2 - 5x + 1$$

$$p_4(x) := 13x^3 - 19x^2 + 8x - 1$$

$$p_5(x) := 13x^3 - 20x^2 + 9x - 1$$

$$p_6(x) := 29x^4 - 59x^3 + 40x^2 - 11x + 1$$

$$p_7(x) := 31x^4 - 61x^3 + 41x^2 - 11x + 1$$

$$p_8(x) := 31x^4 - 63x^3 + 44x^2 - 12x + 1$$

$$p_9(x) := 941x^8 - 3764x^2 + 6349x^6 - 5873x^5$$

$$3243x^4 - 1089x^3 + 216x^2 - 23x + 1$$

We have

Proposition. Let

$$P_{210} := p_0^{67} \cdot p_1^{67} \cdot p_2^{24} \cdot p_3^9 \cdot p_4 \cdot p_5 \cdot p_6^3 \cdot p_7 \cdot p_8 \cdot p_9$$

then

$$(\|P_{210}\|_{[0,1]})^{1/210} = \frac{1}{(2.3543...)}$$

and hence

$$\Omega[0,1] \le \frac{1}{(2.3543\dots)}.$$

Proof. This proof is obviously just a computational verification. It is the algorithm for finding P_{210} which is of some interest. It is based on LLL lattice basis reduction in the following way.

a] Lattice basis reduction finds a short vector in a lattice. If we construct a lattice of the form

$$p(z) \cdot \sum_{k=0}^{n} \alpha_k z^k = \sum_{k=0}^{m} \beta_k z^k$$

where p is a fixed polynomial and the set

$$\{(\alpha_0,\alpha_1,\ldots,\alpha_n)\}$$

is a lattice then the set

$$\{(\beta_0,\beta_1,\ldots,\beta_m)\}$$

is also a lattice, and LLL will return a short vector in the sense of $\sum_{k=0}^{m} |\beta_k|^2$ being relatively small.

Observe that $\left(\sum_{k=0}^{m} |\beta_k|^2\right)^{1/2}$ is just the L_2 norm on the unit disk of the polynomial $\sum_{k=0}^{m} \beta_k z^k$.

So LLL lets us find polynomials of small L_2 norm (and hence small sup norm) on the disk, and we can do this while preserving divisibility by a fixed p.

b] Convert the problem from the interval $[\alpha, \beta]$ to the disk. This is easy. One first maps $[\alpha, \beta]$ to [-2, 2] by a linear change of variables. One then lets x := z + 1/z. This maps a polynomial in x

on [-2,2] to a polynomial in z and 1/z on the boundary of the unit disk.

c] Attack the problem incrementally by using a] and b]. That is, at the k-th stage find a polynomial q_k of degree kN divisible by q_{k-1} of degree (k-1)N using LLL on a lattice of size N+1. This allows us to keep the size of LLL fairly small and uses the fact that integer Chebyshev polynomials tend to have (of necessity) many repeat factors. We used N=10 in the actual computation and started with $q_0:=1$. \square

We can computationally refine the above.

Proposition. The inequality

$$\Omega[0, 1/4] \le \frac{1}{(5.5723\dots)}$$

and hence

$$\Omega[0,1] \le \frac{1}{(2.3605\dots)}.$$

holds.

This is done by minimizing over

$$P_1^{\alpha_1} P_2^{\alpha_2} \dots P_9^{\alpha_9}$$

which is a linear problem.

Corollary. Let k be a positive integer, and let P_{210} be as in the previous Proposition. Then $(P_{210})^k$ divides all the n-th integer Chebyshev polynomials on [0,1] provided n is sufficiently large.

Proof. Each p_i , $i=0,1,\cdots,9$ is irreducible and satisfies

$$p_i(x) = a_k x^k + a^{k-1} x^{k-1} + \dots + a_0$$

with

$$|a_k|^{1/k} < 2.36.$$

Each p_i also has all roots in [0,1]. It follows now by the first Lemma that if Q is a polynomial of degree n with integer coefficients, and

$$\left(\|Q_n\|_{[0,1]}\right)^{1/n} \le \frac{1}{2.3605}$$

then p_i divides Q.

Markov's inequality gives the arbitrarily high multiplicity eventually. \Box

We deduce immediately as above.

Corollary. The polynomials

$$p_0, p_1, \ldots, p_9$$

are the only irreducible polynomials with all their roots in [0,1] of the form

$$p(z) = a_n z^n + a_{n-1} z^{n-1} + \dots + a_0$$

with

$$|a_n|^{1/n} < 2.36.$$

3. Finer Structure.

The exact dependence of $\Omega[a, b]$ on the interval [a, b] is interesting and complicated. If we let

$$\Omega(x) := \Omega[0, x]$$

Then clearly Ω is a non-decreasing function on $(0, \infty)$. Obviously

$$\lim_{x \to 0} \Omega(x) = 0$$

(consider x^m on $[0, \delta]$). So $\Omega(x)$ maps [0, 4] onto [0, 1].

It is an exercise to show that Ω is in fact continuous. This follows mostly from a theorem of Chebyshev that gives

$$||p_n||_{[0,\delta+\epsilon]} \le (1+k_{\epsilon,\delta})^n ||p_n||_{[0,\delta]}$$

for every $p_n \in \mathcal{P}_n$.

What is less obvious is that $\Omega(x)$ is locally flat on many intervals. Indeed it is conceivable that the derivative of Ω is almost everywhere zero.

Theorem. Let $T_n := T_n\{[0,1]\}$ be an n-th integer Chebyshev polynomial on [0,1]. Then T_n is of the form

$$T_n(x) = x^k (1-x)^k S_{n-2k}(x)$$

where (0.26)n < k if n is large enough.

As a consequence, there exists an absolute constant $\delta > 0$ (independent of n) so that T_n is an n-th integer Chebyshev polynomial on larger intervals [-a, 1+a] for every $a \in (0, \delta]$.

4. The Schur-Siegel Trace Problem.

Let $\alpha := \alpha_1$ be an algebraic number with conjugate roots $\alpha_2, \dots, \alpha_n$. We say that α is **totally real (positive)** if all the α_i are real (positive). The **trace** of a totally positive algebraic integer is

$$\alpha_1 + \alpha_2 + \cdots + \alpha_n$$
.

Except for finitely many explicit exceptions, if α is a **totally real algebraic integer** then then

$$\frac{\alpha_1 + \alpha_2 + \dots + \alpha_d}{d} \ge 1.648, \quad \text{Schur (18)}$$

$$\frac{\alpha_1 + \alpha_2 + \dots + \alpha_d}{d} \ge 1.733, \quad \text{Siegel (43)}$$

$$\frac{\alpha_1 + \alpha_2 + \dots + \alpha_d}{d} \ge 1.771, \quad \text{Smyth (83)}.$$

Note that $4\cos^2(\pi/p)$ is a totally positive algebraic integer of degree (p-1)/2 and trace p-2 for p prime. So the best constant in the above theorem is less than 2.

Connection to the integer Chebyshev Problem is

Proposition. If

$$\Omega[0, 1/m] < \frac{1}{m+\delta}.$$

Then, with finitely many exceptions,

$$\frac{\alpha_1 + \alpha_2 + \dots + \alpha_d}{d} \ge \delta$$

for every totally positive algebraic integer α_1 of degree d > 1 with conjugates $\alpha_2, \ldots, \alpha_d$.

Proof. Mostly an application of the original lemma and the Arithmetic-Geometric mean inequality. \Box

Corollary. If α_1 is a totally positive algebraic integer of degree d > 1 with conjugates $\alpha_2, \ldots, \alpha_d$ then

$$\frac{\alpha_1 + \alpha_2 + \dots + \alpha_d}{d} > 1.752$$

with at most finitely many exceptions. (No exceptions of degree greater than 8.)

This is not as good as Smyth's result. It, however, follows immediately from a computation, as in Section 2, which shows that

$$\Omega[0, 1/200] < \frac{1}{201.752}$$

and gives the factors of an example which yields the above upper bound.

5. Open Problems.

There are a myriad of open problems in and around integer Chebyshev polynomials. We formulate a few of them as questions.

- Q1. Find a reasonable algorithm for exactly computing integer Chebyshev polynomials on [0, 1] that would work up to, say, degree 200.
- **Q2.** Are the integer Chebyshev polynomials eventually unique?
- **Q3.** Do the integer Chebyshev polynomials on [0,1] have all their roots in [0,1]?
- **Q4.** Determine $\Omega[0, \alpha]$ exactly for any $0 < \alpha < 4$.

Q5. Determine the limit (or the limsup, if the limit does not exist) of

$$(\Omega[0, 1/m])^{-1} - m.$$

Q6. Are all the irreducible factors of the integer Chebyshev polynomials on [0,1] forced to be factors by as in the first Lemma? That is, are all irreducible factors q of the form

$$q(x) = a_k x^k + a_{k-1} x^{k-1} + \dots + a_0$$

with

$$|a_k|^{1/k} < (\Omega[0,1])^{-1}$$
?

Q7. Show that there exist infinitely many irreducible polynomials with integer coefficients which divide an n-th integer Chebyshev polynomial on [0,1] for some n.

6. Prouhet-Tarry-Escott Problem.

Conjecture.

For any N there exists $p \in Z[x]$ (a polynomial with integer coefficients) so that

$$p(x) = (x-1)^N q(x) = \sum a_k x^k$$

and

$$S(p) := \Sigma |a_k| = 2N.$$

Almost equivalently (though not quite obviously)

$$||p||_{L^2\{|z|=1\}} = \sqrt{2N}.$$

The Basis for the Conjecture.

$$x^{\alpha_1} + \ldots + x^{\alpha_N} - x^{\beta_1} - \ldots - x^{\beta_N} = 0((x-1)^N).$$

For
$$N = 2, \ldots, 10$$
 with

$$[\alpha_1, \ldots, \alpha_N]$$
 and $[\beta_1, \ldots, \beta_N]$

$$[0,3] = [1,2]$$

$$[1, 2, 6] = [0, 4, 5]$$

$$[0, 4, 7, 11] = [1, 2, 9, 10]$$

$$[1, 2, 10, 14, 18] = [0, 4, 8, 16, 17]$$

$$[0, 4, 9, 17, 22, 26] = [1, 2, 12, 14, 24, 25]$$

$$[0, 18, 27, 58, 64, 89, 101]$$

$$= [1, 13, 38, 44, 75, 84, 102]$$

[0, 4, 9, 23, 27, 41, 46, 50]

= [1, 2, 11, 20, 30, 39, 48, 49]

[0, 24, 30, 83, 86, 133, 157, 181, 197]

= [1, 17, 41, 65, 112, 115, 168, 174, 198]

[0, 3083, 3301, 11893, 23314, 24186, 35607,

44199, 44417, 47500] =

[12, 2865, 3519, 11869, 23738, 23762, 35631,

43981, 44635, 47488]

• The size 10 example illustrates the problems inherent with searching for a solution.

Partial History.

- Euler
- Prouhet (1851)
- Tarry (1910) Small Examples
- Escott (1910) Small Examples
- \bullet Letac (1941) Size 9 and 10
- \bullet Gloden (1946) Size 9 and 10
- Smyth (Math Comp. 1991) Size 10 generalized.

Diophantine Form

Find distinct integers $[\alpha_1, \ldots, \alpha_N]$ and $[\beta_1, \ldots, \beta_N]$ so that

$$\alpha_{1} + \ldots + \alpha_{N} = \beta_{1} + \ldots + \beta_{n}$$

$$\alpha_{1}^{2} + \ldots + \alpha_{N}^{2} = \beta_{1}^{2} + \ldots + \beta_{n}^{2}$$

$$\vdots \qquad \vdots \qquad \vdots$$

$$\alpha_{1}^{N-1} + \ldots + \alpha_{N}^{N-1} = \beta_{1}^{N-1} + \ldots + \beta_{N}^{N-1}$$

More Open Questions.

- The problem is completely open for $N \geq 11$.
- We computed extensively on N=11 to show no (symmetric) solutions of degree ≤ 745 .

The Weak Prouhet-Tarry-Escott Problem.

Problem. For fixed N find $p \in Z[x]$

$$p(x) = (x-1)^N q(x) = \sum a_k x^k$$

that minimizes

$$S(p) = \Sigma |a_i|$$

or

$$S^{2}(p) = (\Sigma |a_{i}|^{2})^{1/2}$$

• Solving $S(p) = |S^2(p)|^2 = 2N$ is the Prouhet-Tarry-Escott-Problem and is the big prize.

• Showing that there exist

$${p_N} = {(x-1)^N q(x)}$$

so that

$$S(p_N) = o(N \log N)$$

is also a big prize.

- This shows that the "Easier Waring Problem" is easier than the "Waring Problem" (At the moment.)
- That is: it requires essentially fewer powers to write every integer as sums and differences of Nth powers than just as sums of Nth powers. (Fuchs and Wright, Quart. J. Math. 1936).

• It is known that

$$S((x-1)^N q(x)) \le \frac{N^2}{2}$$

is possible.

Any improvement would be a major step.

• If we demand that p has a zero of order N but not N+1 at 1 then

$$S(p) = 0((\log N)N^2)$$

is possible (Hua).

Any improvement would be interesting.